



DS 102: Data, Inference, and Decisions

Lecture 4

Michael Jordan

University of California, Berkeley

Some Column-Wise Rates

		Decision	
		0	1
Reality	0	n_{00}	n_{01}
	1	n_{10}	n_{11}

$$\text{false discovery proportion} = \frac{n_{01}}{n_{01} + n_{11}}$$

The Goal: Control Errors A Priori

- The row-focused Neyman-Pearson paradigm, with its Type I and Type II errors, provides **a priori control**
 - meaning that if my assumptions about the null and alternative distributions are correct, then I can guarantee that these errors will be small (in an average, frequentist sense---over multiple draws of data)

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- If I'm only testing one hypothesis, that's satisfying
- The problem that arose with our A/B testing example arose because we were doing many tests
- Can we find a way to obtain a priori control when there are many tests?

Comments on the Column-Wise Rates

- They can be thought of as estimates of conditional probabilities
- They **are** dependent on the **prevalence** (i.e., the probabilities of the two states of Reality in the population), via Bayes' Theorem
 - as such, they are more Bayesian
 - this is arguably a good thing
- Notation: let H denote Reality, and let D denote the decision

A Bayesian Calculation

$$P(H = 0 \mid D = 1) = \frac{P(H = 0, D = 1)}{P(D = 1)}$$

A Bayesian Calculation

$$\begin{aligned} P(H = 0 \mid D = 1) &= \frac{P(H = 0, D = 1)}{P(D = 1)} \\ &= \frac{P(D = 1 \mid H = 0)P(H = 0)}{P(D = 1)} \\ &= \frac{P(\text{Type I error}) \cdot \pi_0}{P(D = 1)} \end{aligned}$$

- We could upper bound π_0 with 1, and so the numerator can be controlled; what about the denominator?

A Bayesian Calculation

- Using the law of total probability, we have:

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- So we see that $P(D = 1)$ depends on the prior π_0
- Is this a problem?
 - i.e., do we have to either decide to be Bayesian and supply the prior, or decide to be frequentist and abandon this approach?
- No! Note that it's easy to estimate $P(D = 1)$ directly from the data!

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- This controls the FDR!

P-Values

- Consider a point-null hypothesis, $\theta = 0$, and \mathbb{P} denote that null
- Consider a statistic, $T(X)$, which has a continuous distribution under the null, and let $F(t)$ denote its tail cdf:

$$F(t) = \mathbb{P}(T > t)$$

- Define the P-value as $P = F(T)$
- The P-value has a uniform distribution under the null:

$$\mathbb{P}(P < p) = \mathbb{P}(F(T) < p) = \mathbb{P}(T > F^{-1}(p)) = F(F^{-1}(p)) = p$$

A Generic Decision Rule

- Reject H_i if the random variable T_i is equal to 1:

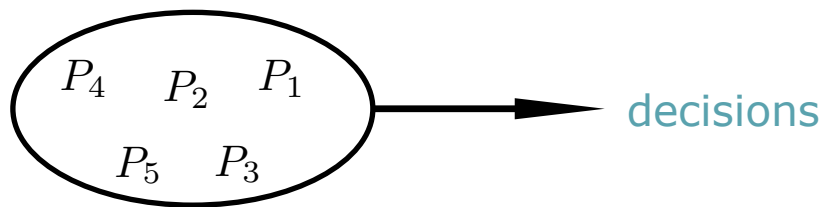
$$T_i = \begin{cases} 1, & \text{if } P_i \leq \alpha_i \\ 0, & \text{otherwise} \end{cases}$$

The Online Problem

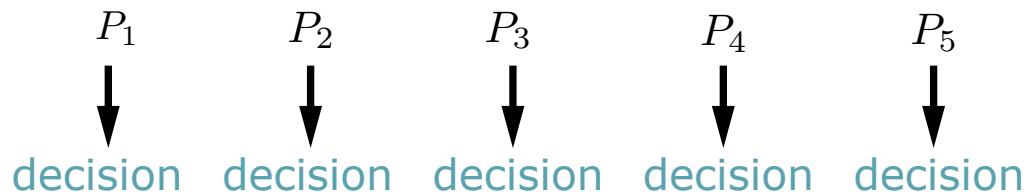
- Classical statistics, and also the Benjamini & Hochberg algorithm focused on a batch setting in which all data has already been collected
- E.g., for Benjamini & Hochberg, you need all of the p-values before you can get started
- Is it possible to consider methods that make sequences of decisions, and provide FDR control at any moment in time
- Is it conceivable that one can achieve **lifetime** FDR control?

Online vs Offline FDR Control

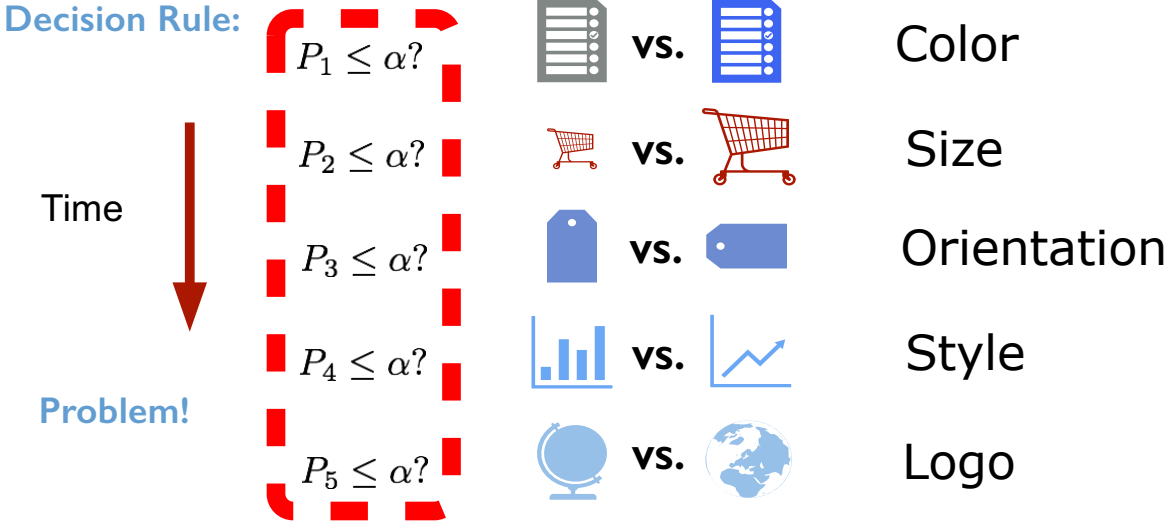
- Classical FDR procedures (such as BH) which make all decisions simultaneously are called “offline”



- “Online” FDR procedures make decisions one at a time



Example: Many Enterprises Run Thousands of So-Called A/B Tests Each Day



Challenges

- It's not clear how to do change batch procedures such as Benjamini-Hochberg procedure to be online

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- It's not clear how to do change batch procedures such as Benjamini-Hochberg procedure to be online
- We might retreat to Bonferroni, which would allow us to set α to $0.05/n$ and thereby have a FWER of 0.05 after n tests
 - but what do we do on the $(n + 1)th$ test?
 - we eventually can't do any more tests
 - we've used up our "alpha wealth"

A More General Approach: Time-Varying Alpha

Decision Rule:

Time



$$P_1 \leq \alpha_1?$$



vs.



Color

$$P_2 \leq \alpha_2?$$



vs.



Size

$$P_3 \leq \alpha_3?$$



vs.



Orientation

$$P_4 \leq \alpha_4?$$



vs.



Style

$$P_5 \leq \alpha_5?$$



vs.



Logo

More Challenges

- We want to keep going for an arbitrary amount of time, so we need $\sum_{t=1}^{\infty} \alpha_t = 1$, and $\sum_{t=1}^T \alpha_t < 1$ for any fixed T
- An example: $\alpha_t = 2^{-t}$
- But now we have less and less power to make discoveries over time, and eventually we may as well quit
- Is there any way out of this dilemma?

A Glimmer of Hope

- Recall that the FDP is a **ratio** of two counts
- We can make a ratio small in one of two ways:
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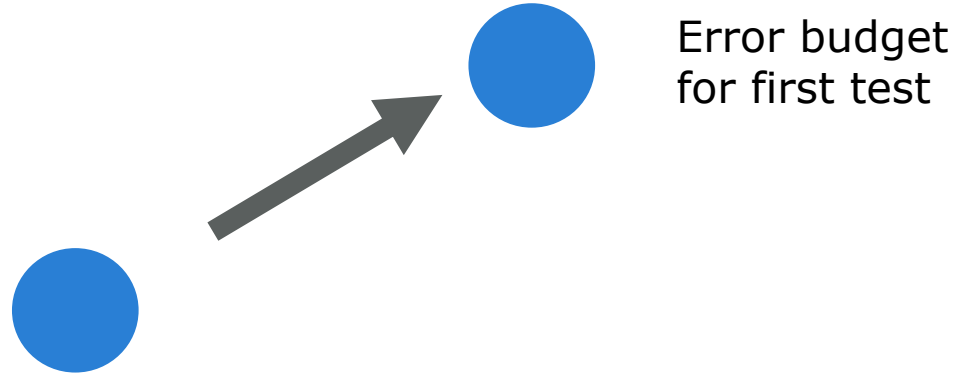
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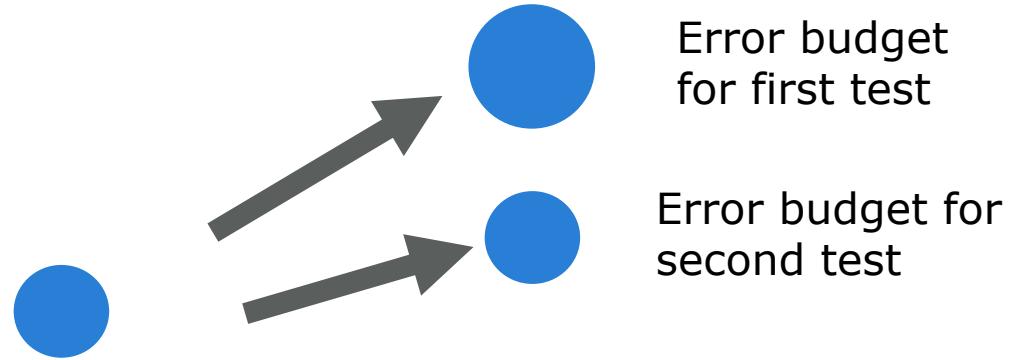
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- The **denominator** can be made large by making lots of discoveries
- Perhaps we can earn a bit of alpha whenever we make a discovery, to be invested and used for false discoveries later

Online FDR Control : High-Level Picture



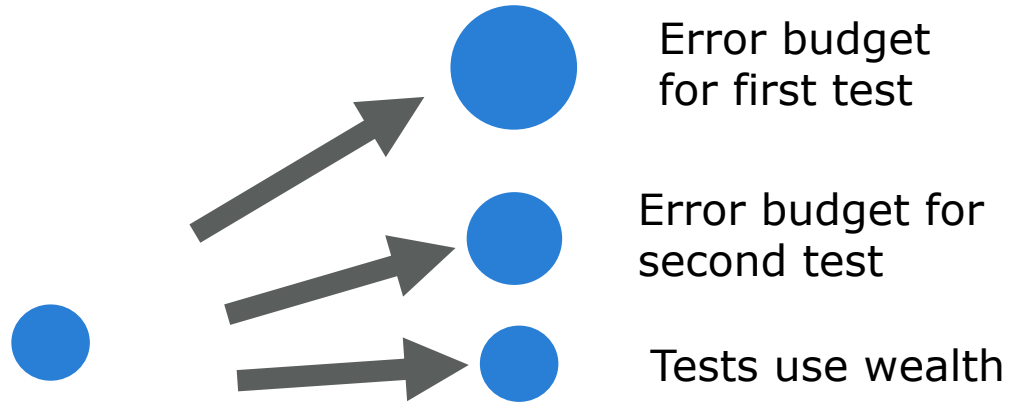
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or "alpha-wealth"

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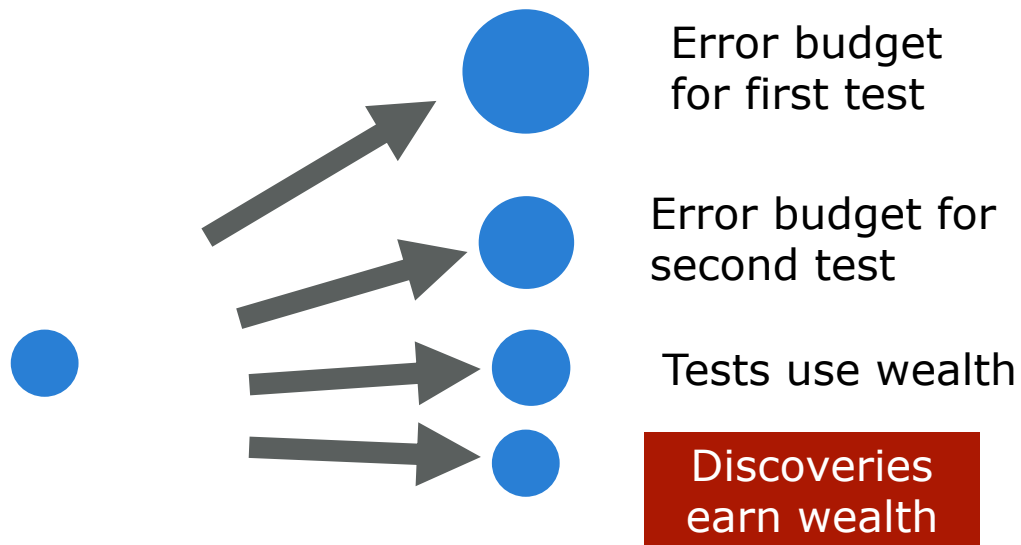
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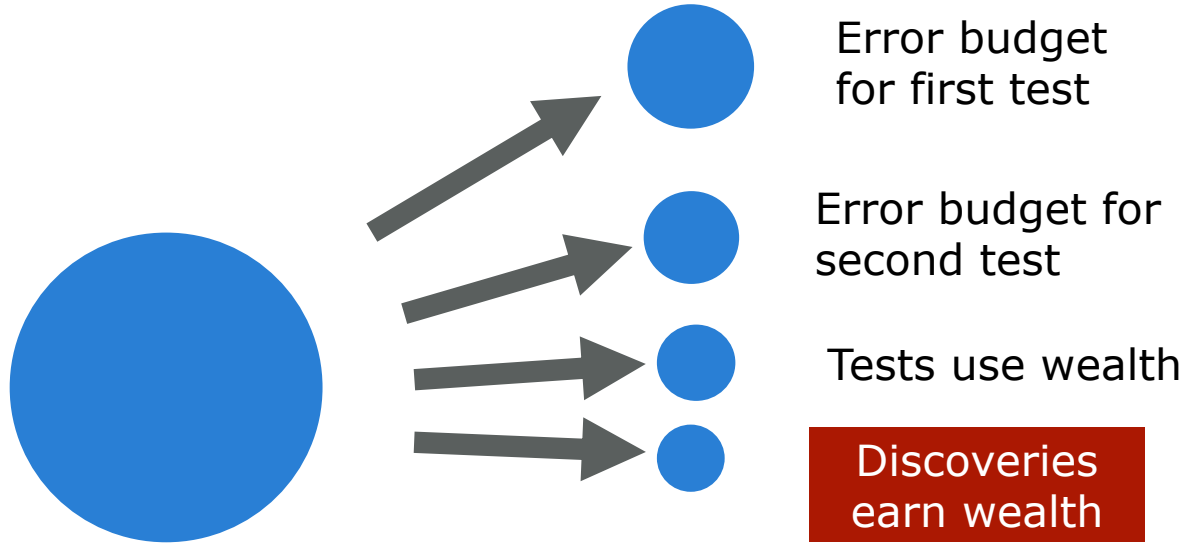
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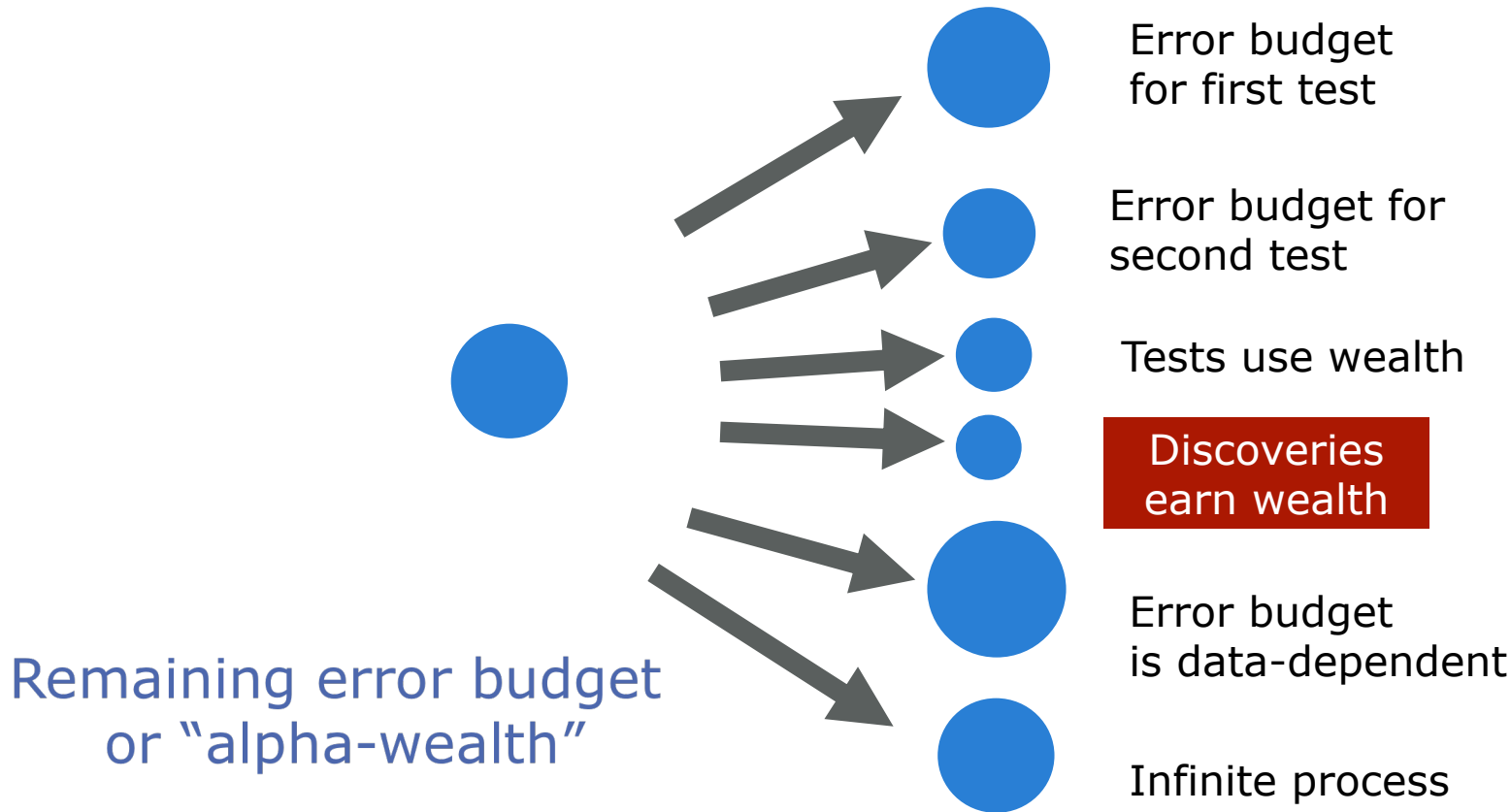
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Online FDR Algorithms

- The first online FDR algorithm was known as “alpha investing” and is due to Foster and Stine (2008)
- A more recent (and simpler) online FDR algorithm is due to Javanmard and Montanari, and is called “LORD”
- The basic idea is to assign α_t in a way that ensures

$$\widehat{\text{FDP}}(t) := \frac{\sum_{i=1}^t \alpha_i}{\sum_{i=1}^t 1\{P_i \leq \alpha_i\}} \leq \alpha$$

Algorithm 1 The LORD Procedure

input: FDR level α , non-increasing sequence $\{\gamma_t\}_{t=1}^{\infty}$ such that $\sum_{t=1}^{\infty} \gamma_t = 1$,
initial wealth $W_0 \leq \alpha$

Set $\alpha_1 = \gamma_1 W_0$

for $t = 1, 2, \dots$ **do**

 p-value P_t arrives

 if $P_t \leq \alpha_t$, reject P_t

$$\alpha_{t+1} = \gamma_{t+1} W_0 + \gamma_{t+1-\tau_1} (\alpha - W_0) \mathbf{1}\{\tau_1 < t\} + \alpha \sum_{j=1}^{\infty} \gamma_{t+1-\tau_j} \mathbf{1}\{\tau_j < t\},$$

 where τ_j is time of j -th rejection $\tau_j = \min\{k : \sum_{l=1}^k \mathbf{1}\{P_l \leq \alpha_l\} = j\}$

end

A Stripped-Down Version of LORD

- Only consider the most recent rejection
- This renews the wealth, which further decays
- See description, and proof of mFDR control, on board

A Heuristic Argument for LORD's Control of FDR

- We make an approximation:

$$\text{FDR} \approx \frac{\mathbb{E}[\sum_{i \leq t, i \text{ null}} 1\{P_i \leq \alpha_i\}]}{\mathbb{E}[\sum_{i \leq t} 1\{P_i \leq \alpha_i\}]}$$

and then compute:

$$\begin{aligned} \mathbb{E} \left[\sum_{i \leq t, i \text{ null}} 1\{P_i \leq \alpha_i\} \right] &= \sum_{i \leq t, i \text{ null}} \mathbb{E}[\mathbb{E}[1\{P_i \leq \alpha_i\} | \alpha_i]] = \sum_{i \leq t, i \text{ null}} \mathbb{E}[\mathbb{P}\{P_i \leq \alpha_i | \alpha_i\}] \\ &= \sum_{i \leq t, i \text{ null}} \mathbb{E}[\alpha_i] \leq \mathbb{E}[\sum_{i \leq t} \alpha_i] \leq \alpha \mathbb{E}[\sum_{i \leq t} 1\{P_i \leq \alpha_i\}] \end{aligned}$$

where the last line uses:

$$\widehat{\text{FDP}}(t) := \frac{\sum_{i=1}^t \alpha_i}{\sum_{i=1}^t 1\{P_i \leq \alpha_i\}} \leq \alpha$$

- This establishes:

$$\text{FDR} \leq \alpha$$